## Distribution on Contingency of Alignment of Two Literal Sequences under Constrains

Running title: Distribution on alignments under constrains

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#### Abstract

The case of ungapped alignment of two literal sequences under constrains is considered. The analysis lead to general formulas for probability mass function and cumulative distribution function for the general case of using an alphabet with a chosen number of letters (e.g. 4 for deoxyribonucleic acid sequences) in the expression of the literal sequences. Formulas for three statistics including mean, mode, and standard deviation were obtained. Distributions are depicted for three important particular cases: alignment on binary sequences, alignment of trinomial series (such as coming from generalized Kronecker delta), and alignment of genetic sequences (with four


literals in the alphabet). A particular case when sequences contain each letter of the alphabet at least once in both sequences has also been analyzed and some statistics for this restricted case are given.

Keywords: alignment; contingency matrix; probability mass function (PMF); cumulative distribution function (CDF).

## 1 Introduction

Researches related to sequence alignments are frequently done due to the huge amount of already identified sequence of DNA (deoxyribonucleic acid), RNA (ribonucleic acid), or proteins (Pruitt et al. 2012). Sequence alignment is defined as a way of arrange DNA, RNA (Allali et al. 2012), or amino acid (Mongiovì and Sharan 2013) sequences to identify similar regions that could reflect functional, structural or evolutionary relationships between sequences (Mount 2004). Several algorithms were developed and implemented for global (Rahrig et al. 2013; Szalkowski and Anisimova 2013) or local alignments (Phuong et al. 2006; Tabei and Asai 2009; Frith et al. 2010), each algorithm with certain advantages and disadvantages. For example, the approach proposed by Szalkowski and Anisimova (2013) detect insertions and deletions of TR (tandem repeats) units not restricted to TR unit boundaries and proved more performing ( $\sim 10 \%$ ) compared to other aligners for cases with divergence high TR rates. Frith et al. (2010) assessed several combinations of score parameters for alignment (495) and found that high value of X-drop parameter are not always better, when tandem repeats are masked in a non-standard way the E-values accurately indicate the rate of spurious alignment, while highly reliable subsets of aligned bases could be obtained by $\gamma$ centroid alignment.

The state-of-the art in the pairwise alignments showed that statistical significance of the alignment is directly related to scoring scheme, sequence length and number of literals in the sequence (Mott 2005). Different approaches had been developed and implemented to estimate statistical significance of scores of alignment. BLAST2.0 (Altschul et al. 1997) implement a lookup method where $K$ and $\lambda$ parameters are pre-computed for different scoring schemes using average amino acid composition of both sequences. FASTA package (PRSS program) estimate the statistical significance of the shuffled (1,000 times) score distribution (Pearson 2000) while HMMER use maximum likelihood fitting in estimation of statistical significance (Mitrophanov and Borodovsky 2006).

Different scoring functions, such as probability consistency transformation - PCT (Do et al. 2005), Burrows-Wheeler Transform - BWT (Li and Durbin 2009), PSAR (Kim and Ma 2011), gPSP (Mokaddeml and Elloumi 2013), PSAR-Align (Kim and Ma 2014), etc., are used to characterize the alignment. BWT could be seen as a scoring function since the extension of BWT introduced by Mantaci et al. (2008) led to a general method for comparing sequences.

Distribution analysis found its usefulness in assessment of different natural (Bolboacă et al. 2011, Jäntschi et al. 2012a; 2011) or simulated phenomena (Jäntschi and Bolboacă 2011; Jäntschi et al. 2012b). Karlin and Altschul (1990) proved for ungapped alignment that the optimal local alignment scores used in the evaluation of sequence alignments follow an extreme-value distribution (characterized by characteristic value $K$, and scale $\lambda$ ). Computational experiments suggest that the optimal local alignment scores also apply to gapped local alignments (Smith et al. 1985; Altschul et al. 2001). The most important advantages of the island method (Olsen et al. 1999) over the direct method (Waterman 1994) in estimation of statistical parameters for gapped local sequence alignment is related with systematic errors that are easiest to be controlled (Altschul et al. 2001).

Our research started from the hypothesis that the distribution of ungapped alignments could provide useful information about the chance of their occurrences. A statistical approach based on distribution analysis able to identify the thresholds for rejecting an ungapped alignment by chance has been developed and is presented in this manuscript.

## 2 Methods

The global ungapped alignment of two sequences of equal length (equal number of literals, $n$ ) was investigated in this study. The general formulation of the problem investigated in this study along with an example of a particular case (total number of letters $q=4-$ for DNA and RNA, A $=$ adenine, $\mathrm{C}=$ cytosine, $\mathrm{G}=$ guanine, and $\mathrm{U}=$ uracil; respectively $\mathrm{A}, \mathrm{C}, \mathrm{G}$, and $\mathrm{T}=$ thymine) is presented in Table 1.

Table 1: Alignment of two sequences of identical length: general case (left-hand) and particular case (right-hand, $q=4$ )

| $j$ | 1 | 2 | 3 | 4 | $\ldots$ | $n-4$ | $n-3$ | $n-2$ | $n-1$ | $n$ |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Seq1 | $\mathrm{a}_{11}$ | $\mathrm{a}_{12}$ | $\mathrm{a}_{13}$ | $\mathrm{a}_{14}$ | $\ldots$ | $\mathrm{a}_{1(\mathrm{n}-4)}$ | $\mathrm{a}_{1(\mathrm{n}-3)}$ | $\mathrm{a}_{1(\mathrm{n}-2)}$ | $\mathrm{a}_{1(\mathrm{n}-1)}$ | $\mathrm{a}_{\mathrm{n}}$ |
| Seq2 | $\mathrm{a}_{21}$ | $\mathrm{a}_{22}$ | $\mathrm{a}_{23}$ | $\mathrm{a}_{24}$ | $\ldots$ | $\mathrm{a}_{2(\mathrm{n}-4)}$ | $\mathrm{a}_{2(\mathrm{n}-3)}$ | $\mathrm{a}_{2(\mathrm{n}-2)}$ | $\mathrm{a}_{2(\mathrm{n}-1)}$ | $\mathrm{a}_{\mathrm{n}}$ |
| Match | $?$ | $?$ | $?$ | $?$ | $\ldots$ | $?$ | $?$ | $?$ | $?$ | $?$ |

Seq1 $=$ first sequence; Seq2 $=$ second sequence;

| $j$ | 1 | 2 | 3 | 4 | $\ldots$ | $n-4$ | $n-3$ | $n-2$ | $n-1$ | $n$ |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Seq1 | A | C | C | G | $\ldots$ | U | A | G | A | C |
| Seq2 | C | A | C | U | $\ldots$ | A | A | G | C | A |
| Match | no | no | yes | no | $\ldots$ | no | yes | yes | no | no |

$\mathrm{a}_{\mathrm{ij}}$ : the first subscript number refers the number of sequence (1 or 2 ); the second subscript number refers the index of the literal in the sequence $(1 \leq j \leq n)$;
$n$ = the length of the sequence

A match is present when identical alphabet letters are present in both sequences at the same. The case presented in Table 1 could be transposed in a contingency of alignment as it is presented in Table 2. The total number of possible literals in the sequences (q) gives the size of alignment contingency (Table 2).

Table 2: Alignment contingency: general case (left-hand) and particular case (right-hand, $q=4$ )


The number of aligned literals from two stings of $n$ literals, $P S q(q)$, is therefore given by the main diagonal of the alignment contingency presented in Table 2 (Eq(1) for general case):

$$
\begin{equation*}
\operatorname{PSq}(q)=\Sigma_{1 \leq i \leq n} b_{11}+\Sigma_{1 \leq i \leq n} b_{22}+\ldots+\Sigma_{1 \leq i \leq n} b_{(q-1)(q-1)}+\Sigma_{1 \leq i \leq n} b_{q q} \tag{1}
\end{equation*}
$$

where $P S q(q)$ ranges from 0 (no matches) to $n$ (perfect alignment).
In the particular case ( $q=4$ ), the $\mathrm{Eq}(1)$ became:

$$
\begin{equation*}
\operatorname{PSq}(4)=\Sigma_{v \in\{A, C, G, U\}} \Sigma_{1 \leq i \leq n}\left(\text { Seq } 1_{i}=v, \operatorname{Seq} 2_{i}=v\right) \tag{1P}
\end{equation*}
$$

Based on the perfect alignment (the same literal exists at the same position in both sequences), the alignment ratio (AR) could be obtained for the general case using the formulas presented in $\mathrm{Eq}(2)$ :

$$
\begin{equation*}
A R=P S q(q) / n \tag{2}
\end{equation*}
$$

and for the particular case ( $q=4$ ) using the formula presented in $\mathrm{Eq}(2 \mathrm{P})$ :

$$
\begin{equation*}
A R=P S q(4) / n \tag{2P}
\end{equation*}
$$

Two cases of alignment of two sequences of equal length were investigated in this study:

- Unrestricted case: no restriction is imposed in regards of appearance of letters in each sequence.
- Restricted case: all letters of the alphabet (genetic sequence alignment: A, C, G, T or A, C, G, U) appear at least once in each of both sequences.

A full enumeration study was conducted on ungapped alignment of two sequences with identical length from 2 to 10 , with 2 to 4 literals in the alphabet ( $0 \leq i \leq 10$, where $i=$ number of matches). The full enumeration study was conducted in accordance with conventional procedures to generate all numbers whose representation in base $q$ (number of literals in the alphabet, $2 \leq n \leq 4$ ) has $n$ (the length of the sequence) digits; two by two such sequences were generated and then aligned. All frequencies of alignments in the generated sequences were counted for unrestricted case and for restricted case (where were counted only if accomplished the criterion of all letters appearance).

The results of the full enumeration analysis were used to identify (whenever possible) the general formulas for:

- Number of possibilities of arrangements $S q(i ; n, q)$;
- Number of perfect alignments $S q(i=n ; n, q)$ and number of no matches $S q(i=0 ; n, q)$;
- Statistical parameters of alignment: mode, mean, and variance;
- Probability mass function (PMF) and cumulative distribution function (CDF);
- Thresholds for alignment by chance ( $q=4$ and $4 \leq n \leq 40$ ) at a significance level of $5 \%$ by MonteCarlo experiment.


## 3 Results and Discussion

The results and associated discussion presented in this section refer to the unrestricted case as well as, in certain case, to the restricted case (the restriction referred to apparition at least once of any letters from alphabet in each sequence, equations referred with ' R ' along the manuscript). The results obtained on full enumeration study are given in Tables 3-5.

Table 3: Total number of possible arrangements: $2 \leq q \leq 4$ and $n \leq 10$

| n | Unrestricted case |  |  | Restricted case |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $q=2$ | $q=3$ | $q=4$ | $q=2$ | $\boldsymbol{q}=3$ | $q=4$ |
| 2 | 16 |  |  | 4 |  |  |
| 3 | 64 | 729 |  | 36 | 36 |  |
| 4 | 256 | 6561 | 65536 | 196 | 1296 | 576 |
| 5 | 1024 | 59049 | 1048576 | 900 | 22500 | 57600 |
| 6 | 4096 | 531441 | 16777216 | 3844 | 291600 | 2433600 |
| 7 | 16384 | 4782969 | 268435456 | 15876 | 3261636 | 70560000 |
| 8 | 65536 | 43046721 | 4294967296 | 64516 | 33593616 | 1666598976 |
| 9 | 262144 | 387420489 | 68719476736 | 260100 | 329422500 | 34774790400 |
| 10 | 1048576 | 3486784401 | 1099511627776 | 1044484 | 3133760400 | 669974990400 |

Table 4: Full enumeration results on unrestricted case: $q=2$ and $q=4, q \leq n \leq 10$

| $q=2$ and $2 \leq n \leq 10$ |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| n | $i=0$ | $i=1$ | $i=2$ | $i=3$ | $i=4$ | $i=5$ | $i=6$ | $i=7$ | $\boldsymbol{i}=8$ | $i=9$ | $i=10$ |
| 2 | 4 | 8 | 4 |  |  |  |  |  |  |  |  |
| 3 | 8 | 24 | 24 | 8 |  |  |  |  |  |  |  |
| 4 | 16 | 64 | 96 | 64 | 16 |  |  |  |  |  |  |
| 5 | 32 | 160 | 320 | 320 | 160 | 32 |  |  |  |  |  |
| 6 | 64 | 384 | 960 | 1280 | 960 | 384 | 64 |  |  |  |  |
| 7 | 128 | 896 | 2688 | 4480 | 4480 | 2688 | 896 | 128 |  |  |  |
| 8 | 256 | 2048 | 7168 | 14336 | 17920 | 14336 | 7168 | 2048 | 256 |  |  |
| 9 | 512 | 4608 | 18432 | 43008 | 64512 | 64512 | 43008 | 18432 | 4608 | 512 |  |
| 10 | 1024 | 10240 | 46080 | 122880 | 215040 | 258048 | 215040 | 122880 | 46080 | 10240 | 1024 |
| $q=4$ and $4 \leq n \leq 10$ |  |  |  |  |  |  |  |  |  |  |  |
| 4 | 20736 | 27648 | 13824 | 3072 | 256 |  |  |  |  |  |  |
| 5 | 248832 | 414720 | 276480 | 92160 | 15360 | 1024 |  |  |  |  |  |
| 6 | 2985984 | 5971968 | 4976640 | 2211840 | 552960 | 73728 | 4096 |  |  |  |  |
| 7 | 35831808 | 83607552 | 83607552 | 46448640 | 15482880 | 3096576 | 344064 | 16384 |  |  |  |
| 8 | 429981696 | 1146617856 | 1337720832 | 891813888 | 371589120 | 99090432 | 16515072 | 1572864 | 65536 |  |  |
| 9 | 5159780352 | 15479341056 | 20639121408 | 16052649984 | 8026324992 | 2675441664 | 594542592 | 84934656 | 7077888 | 262144 |  |
| 10 | 61917364224 | 206391214080 | 309586821120 | 275188285440 | 160526499840 | 64210599936 | 17836277760 | 3397386240 | 424673280 | 31457280 | 1048576 |

Table 5: Full enumeration results on restricted case: $q=2$ and $q=3, q \leq n \leq 10$

| $q=2$ |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| n | $i=0$ | $\boldsymbol{i}=1$ | $\boldsymbol{i}=2$ | $\boldsymbol{i}=3$ | $i=4$ | $i=5$ | $i=6$ | $\boldsymbol{i}=7$ | $\boldsymbol{i}=8$ | $\boldsymbol{i}=9$ | $i=10$ |
| 2 | 2 | 0 | 2 |  |  |  |  |  |  |  |  |
| 3 | 6 | 12 | 12 | 6 |  |  |  |  |  |  |  |
| 4 | 14 | 48 | 72 | 48 | 14 |  |  |  |  |  |  |
| 5 | 30 | 140 | 280 | 280 | 140 | 30 |  |  |  |  |  |
| 6 | 62 | 360 | 900 | 1200 | 900 | 360 | 62 |  |  |  |  |
| 7 | 126 | 868 | 2604 | 4340 | 4340 | 2604 | 868 | 126 |  |  |  |
| 8 | 254 | 2016 | 7056 | 14112 | 17640 | 14112 | 7056 | 2016 | 254 |  |  |
| 9 | 510 | 4572 | 18288 | 42672 | 64008 | 64008 | 42672 | 18288 | 4572 | 510 |  |
| 10 | 1022 | 10200 | 45900 | 122400 | 214200 | 257040 | 214200 | 122400 | 45900 | 10200 | 1022 |
| $\boldsymbol{q}=3$ |  |  |  |  |  |  |  |  |  |  |  |
| 3 | 12 | 18 | 0 | 6 |  |  |  |  |  |  |  |
| 4 | 288 | 504 | 324 | 144 | 36 |  |  |  |  |  |  |
| 5 | 3180 | 7410 | 7020 | 3660 | 1080 | 150 |  |  |  |  |  |
| 6 | 26640 | 77220 | 94230 | 63000 | 24570 | 5400 | 540 |  |  |  |  |
| 7 | 195132 | 671622 | 996030 | 828030 | 417690 | 128646 | 22680 | 1806 |  |  |  |
| 8 | 1326528 | 5262768 | 9159192 | 9139536 | 5721660 | 2305296 | 586152 | 86688 | 5796 |  |  |
| 9 | 8624460 | 38653578 | 77098392 | 89828424 | 67375476 | 33753132 | 11308248 | 2449656 | 312984 | 18150 |  |
| 10 | 54532080 | 272115900 | 611443890 | 814654800 | 712684980 | 427812840 | 178525620 | 51181200 | 9664110 | 1089000 | 55980 |

The general formula for the total number of possibilities of arrangements between two sequences of identical length $(S q(i ; n, q))$ in unrestricted case was identified and is given by $\mathrm{Eq}(3)$ :

$$
\begin{equation*}
\sum_{\mathrm{i}=0}^{\mathrm{n}} \mathrm{Sq}(\mathrm{i} ; \mathrm{n}, \mathrm{q})=\left(\mathrm{q}^{\mathrm{n}}\right)^{2} \tag{3}
\end{equation*}
$$

where $i=$ number of matches, $n=$ length of sequence, $q=$ number of letters in the alphabet.

The total number of possibilities of arrangements between two sequences of identical length $(n)$ in restricted case is given by $\operatorname{Eq}(3 \mathrm{R})$ along with formula for $\mathrm{q}=2$ :

$$
\sum_{i=0}^{n} S q(i ; n, q)=\left(\sum_{k=0}^{q-1}(-1)^{k} \cdot\binom{q}{k} \cdot(q-k)^{n}\right)^{2}=\left(q!\cdot\left\{\begin{array}{l}
n  \tag{3R}\\
q
\end{array}\right\}\right)^{2}
$$

where $i=$ number of matches, $n=$ length of sequences, $q=$ number of letters in the alphabet, $k=$ integer, $\binom{\mathrm{q}}{\mathrm{k}}=$ the number of $q$-combinations from a given set $k$ of $n(\mathrm{k}!/((\mathrm{k}-\mathrm{q})!\cdot \mathrm{q}!))$, and $\left\{\begin{array}{l}n \\ \mathrm{q}\end{array}\right\}=$ the number of ways to partition a set of n objects into k non-empty subsets, or Stirling number of the second kind (Sharp 1968).

The formulas on three restricted particular cases of $\mathrm{Eq}(3 \mathrm{R})$ for $2 \leq q \leq 4$ are given bellow:

$$
\sum_{i=0}^{n} S q(i ; n, q)=\left\{\begin{array}{ccc}
\left(q^{n}-q\right)^{2} & \text { for } & q=2 \\
\left(q^{n}-q \cdot 2^{n}+q\right)^{2} & \text { for } & q=3 \\
\left(q^{n}-q \cdot 3^{n}+(q+2) \cdot 2^{n}-q\right)^{2} & \text { for } & q=4
\end{array}\right.
$$

The formula is more complex for the restricted case compared to unrestricted case, as could be observed when $\mathrm{Eq}(3)$ is compared with $\mathrm{Eq}(3 \mathrm{R})$. As expected, the total number of possible arrangements is smaller for restricted case compared to unrestricted case and the difference increased with $q$ (total number of letters in the alphabet). Furthermore, the difference for the same $q$ decreases with the increasing of sample size ( $n$ ), while the decrease is faded for large $n$.

Formula presented in Eq(3R) could be checked for given $n$ and $q$ with the full enumeration results presented in Table 3.

General formula for the total number of $i$ matches out of $n$ for unrestricted case $(\mathrm{Eq}(4))$ and some particular formulas for restricted case ( $2 \leq q \leq 4$ ) (coming from $\left.q^{2 n}=q^{n} \cdot((q-1)+1)^{n}, E q(4 R)\right)$ were obtained as follows:

$$
\begin{equation*}
\operatorname{Sq}(\mathrm{i} ; \mathrm{n}, \mathrm{q})=\mathrm{q}^{\mathrm{n}} \cdot(\mathrm{q}-1)^{\mathrm{n}-\mathrm{i}} \cdot 1^{\mathrm{i}} \cdot\binom{\mathrm{n}}{\mathrm{i}}=\mathrm{q}^{\mathrm{n}} \cdot \operatorname{NewtonBin}(\mathrm{i}, \mathrm{q}-1,1, \mathrm{n}) \text {, where } \tag{4}
\end{equation*}
$$

$\operatorname{NewtonBin}(i, x, y, n)=\binom{n}{i} \cdot x^{n-i} \cdot y^{i}$

$$
\begin{align*}
& \operatorname{Sq}(\mathrm{i} ; \mathrm{n}, \mathrm{q}=2)=\left\{\begin{array}{l}
\mathrm{q}^{\mathrm{n}}-\mathrm{q}, \mathrm{i}=0 \\
\left(\mathrm{q}^{\mathrm{n}}-\mathrm{q}^{2}\right) \cdot\binom{\mathrm{n}}{\mathrm{i}}, 0<\mathrm{i}<\mathrm{n} \\
\mathrm{q}^{\mathrm{n}}-\mathrm{q}, \mathrm{i}=\mathrm{n}
\end{array}\right. \\
& \mathrm{Sq}(\mathrm{i} ; \mathrm{n}, \mathrm{q}=3)=\left\{\begin{array}{l}
(2 \cdot \mathrm{q})^{\mathrm{n}}-2 \cdot \mathrm{q} \cdot 4^{\mathrm{n}}+2 \cdot \mathrm{q} \cdot 3^{\mathrm{n}}+\mathrm{q} \cdot 2^{\mathrm{n}}-2 \cdot \mathrm{q}, \mathrm{i}=0 \\
\text { to be determined, } 0<\mathrm{i}<\mathrm{n} \\
\mathrm{q}^{\mathrm{n}}-\mathrm{q} \cdot 2^{\mathrm{n}}+\mathrm{q}, \mathrm{i}=\mathrm{n}
\end{array}\right. \\
& \text { Sq(i;n,q=4)=\{\{} \begin{array}{l}
\text { to be determined, } \mathrm{i}=0 \\
\text { to be determined, } 0<\mathrm{i}<\mathrm{n} \\
\mathrm{q}^{\mathrm{n}}-\mathrm{q} \cdot 3^{\mathrm{n}}+(\mathrm{q}+2) \cdot 2^{\mathrm{n}}-\mathrm{q}, \mathrm{i}=\mathrm{n}
\end{array} \tag{4R}
\end{align*}
$$

where $i=$ number of matches, $q=$ number of letters in the alphabet, $n=$ length of the sequences. The number of perfect matches $(i=n)$ in both unrestricted and restricted case is the square root of the total number of possibilities of arrangements as could be verified with the results presented in Table 4 and 5:

$$
\operatorname{Sq}(\mathrm{i}=\mathrm{n} ; \mathrm{n}, \mathrm{q})=\sqrt{\sum_{\mathrm{i}=0}^{\mathrm{n}} \mathrm{Sq}(\mathrm{i} ; \mathrm{n}, \mathrm{q})}
$$

It should be noted that (see 5\&5R) the problem of counting the number of all full alignments (when $\mathrm{i}=\mathrm{n}$ ) is equivalent with the problem of counting the number of partitions of an n-set into k nonempty but distinguishable boxes (ordered non-empty subsets). The explicit formula for them can be obtained by applying the principle of inclusion-exclusion, when using the universal set consisting of all partitions of the n -set into k (possibly empty) distinguishable boxes, and the exclusion property that the partition has the associated box empty, the principle of inclusion-exclusion gives the answer for the related result (Brualdi 2010). Each arrangement (from the ones of which number is given by 5\&5R) can be seen as an individual sequence obeying the imposed rule to contain all literals and here is a bijective function (if exists a sequence containing all letters and following the imposed rule then exists a perfect arrangement having the second sequence identical with the first one and vice versa). Therefore, the number of paired sequences is the square of this number and the proof for the (3R) formula is completed too.

One important property that can be observed by analyzing the results presented in Table 4 is that, for $q=2$ and $0 \leq i \leq 10$, the distribution of total number of matches is symmetric and had one pick for even $n$ and two equal values at the pick of the distribution for odd $n$. The distribution of total number of matches become asymmetrical and always have just one pick for $q=4$ and $0 \leq i \leq 10$. The second property that could be observed by analyzing the results presented in Table 4 refers the total number of matches for $i=n$ which verify the formula $q^{n}$ for both $q=2$ and $q=4$. When two letters are in the alphabet ( $\mathrm{q}=2$ ), the symmetry of the distribution associated to the total number of matches observed for unrestricted case (Table 4) has also been observed for the restricted case (Table 5). Moreover, in the restricted case for $q=2$, one pick is observed for even $n$ and two equal values at the pick of the distribution for odd $n$. Similar with the non-restricted case, the distribution of the total number of matches become asymmetrical for $q=3$ and with just one value to the pick of the distribution and with a tendency to symmetry for large sample sizes.

The probability mass function (PMF) associated with the frequency of apparition of matches using freely $q$ literals (unrestricted case) is given by Eq(6) while the cumulative distribution function (CDF) is given by $\mathrm{Eq}(7)$.

$$
\begin{align*}
& \operatorname{PMF}_{S q(i, n, q)}=\frac{q^{n} \cdot(q-1)^{n-i} \cdot 1^{i}}{q^{2 n}} \cdot\binom{n}{i}=\frac{(q-1)^{n-i}}{q^{n}} \cdot\binom{n}{i}  \tag{6}\\
& \operatorname{CDF}_{S q(i ; n, q)}=\sum_{k=0}^{i} \operatorname{PDF}_{S q(i, n, q)}=\sum_{k=0}^{i} \frac{(q-1)^{n-k}}{q^{n}} \cdot\binom{n}{k}=\frac{(q-1)^{n}}{q^{n}} \sum_{k=0}^{i}(q-1)^{-k} \cdot\binom{n}{k} \tag{7}
\end{align*}
$$

Graphical representation of PMF for $2 \leq q \leq 4$ is showed in Figure 1 (Eq(6)), while Figure 2 showed the $\operatorname{CDF}(\operatorname{Eq}(7))$.

## [Place Figure 1 here]

Figure 1. Probability mass function for $2 \leq q \leq 4$ (0.0-0.1-0.2-0.3-0.4-0.5 with red, green, blue, cyan and magenta) with PMF on vertical axis and number of aligned literals (i) on frontal axis
(unrestricted case)
[Place Figure 2 here]

Figure 2. Cumulative distribution function for $2 \leq q \leq 4$ (0.0-0.1-0.2-0.3-0.4-0.5 with red, green, blue, cyan and magenta) with CDF on vertical axis and number of aligned literals (i) on frontal axis (unrestricted case)

Figure 1 shows the symmetrical distribution for the unrestricted case as already observed and the remoteness from the symmetry with the increase of the number of letters in the alphabet. The PMF and CDF associated with the frequency of apparition of matches for restricted case is under investigation in our laboratory and the absence of the results is due to our available power of computation.

General formulas of three alignment statistical parameters (named mode, mean, and variance) associated to $\operatorname{Sq}(i ; n, q)$ have been identified for the unrestricted case and are given by $\operatorname{Eq}(8)$ Eq(10).

$$
\begin{align*}
& \hat{\mu}=\left\{\begin{array}{l}
\{\mathrm{k}, \mathrm{k}+1\}, \mathrm{n}=\mathrm{q} \cdot \mathrm{k}+\mathrm{q}-1 \\
\mathrm{k}, \mathrm{n} \neq \mathrm{q} \cdot \mathrm{k}+\mathrm{q}-1
\end{array}\right.  \tag{8}\\
& \mu=\frac{\mathrm{n}}{\mathrm{q}}  \tag{9}\\
& \sigma^{2}=\frac{\mathrm{n} \cdot(\mathrm{q}-1)}{\mathrm{q}^{2}} \tag{10}
\end{align*}
$$

where $\hat{\mu}=$ mode, $\mu=$ arithmetic mean, $q=$ number of letters in the alphabet, $n=$ length of the sequences, $k=$ integer, $\sigma^{2}=$ variance.

Formulas presented in $\operatorname{Eq(8)-Eq(10)~become~the~well~known~binominal~formulas~for~} q=1 / \mathrm{p}(\mathrm{p} \in[0$, 1], probability of success). Therefore, the unrestricted case could be seen as a binomial experiment. The mode - Eq(8) - defines the alignment with highest probability to be observed by chance. The distribution of the mode in unrestricted case for $2 \leq q \leq 4$ and $q \leq n \leq 10$ is presented in Figure 3.

## [Place Figure 3 here]

Figure 3. Mode in distribution of the alignment by chance vs. length of sequences: unrestricted case

Figure 3 showed different number of modes for different length of the sequence and $2 \leq q \leq 4$. The alignment by chance proved systematically bimodal to every odd $n$ when $q=2(n=q \cdot k+(q-1), k$ being any positive integer). For $q=3$, first bimodal alignment appeared when $n=5$ and occurred with a step equal with $q(n=q \cdot k+(q-1))$. For $q=4$, first bimodal alignment appeared when $n=7$ and occurred with a step equal with $q(n=q \cdot k+(q-1))$.

In the restricted case, the formula for the mean proved the same as for unrestricted while the mode in the distribution of alignment by chance is:
$\div q=2$ : bimodal distribution for odd $n$ at $(n-1) / 2$ and $(n+1) / 2$ and unimodal for even $n$ (excepting $n=2$ ) at $n / 2$.
$\div \quad q=3$ : unimodal distribution when $n$ ranges from $q \cdot k$ to $q \cdot k+q-1$ at $k$.
$\div q=4$ (and higher): the expression of mode for the distribution of alignment by chance has not yet been identified.

The threshold of the alignment by chance $\mathrm{CDF}_{95}\left(k, \mathrm{CDF}_{\mathrm{Sq}(\mathrm{k} ; \mathrm{n}, \mathrm{q}) \geq 0.95}\right)$ ) for restricted case when $2 \leq q \leq 4$ had been estimated from a Monte-Carlo experiment and is approximated by $\mathrm{Eq}(11)-\mathrm{Eq}(13)$.

$$
\begin{align*}
& q=2: \mathrm{k}_{\mathrm{CDF} 295 \%}=(0.84 \pm 0.05) \cdot(1+\mathrm{n})^{0.91 \pm 0.02} ; k \rightarrow 0.56 \cdot n \text {, where } n \rightarrow \infty  \tag{11}\\
& q=3: \mathrm{k}_{\mathrm{CDF} 295 \%}=(0.68 \pm 0.05) \cdot(1+\mathrm{n})^{0.89 \pm 0.02} ; k \rightarrow 0.39 \cdot n \text {, where } n \rightarrow \infty  \tag{12}\\
& q=4: \mathrm{k}_{\mathrm{CDF} 295 \%}=(0.66 \pm 0.07) \cdot(1+\mathrm{n})^{0.83 \pm 0.03} ; k \rightarrow 0.30 \cdot n \text {, where } n \rightarrow \infty \tag{13}
\end{align*}
$$

where $q$ = number of letters in the alphabet, $n=$ length of the sequences.
The formulas presented in $\mathrm{Eq}(11)-\mathrm{Eq}(13)$ provide the threshold $\left(\mathrm{CDF}_{95}\right)$ at which, with a risk smaller than the significance level (in this case a significance level of $5 \%$ was used), the obtained matches did not appear by chance.

The distribution of $\mathrm{CDF}_{95}$ for restricted case, $q=4$ and $q \leq n \leq 40$ obtained through simulations is presented in Figure 4.
[Place Figure 4 here]
Figure 4. Thresholds to reject matches by chance for two equally length sequences and $q=4$

The plot presented in Figure 4 showed that for example when $4 \leq n \leq 6$ if there are observed more than three matches, with a $5 \%$ risk to be in error, these matches are not by chance. With two exceptions (when $10 \leq n \leq 11$ and $30 \leq n \leq 33$ ), the thresholds to reject matches by chance in ungapped alignment of two equally length sequences repeatedly are the same for three sample sizes (e.g. $4 \leq n \leq 6$ ).

The genetic sequence responsible for alpha hemoglobin stabilizing protein (AHSP) on Homo sapiens (HS, chromosome 16) and Mustela putorius furo (MP) were used to exemplify the usefulness of CDF on restricted case. One hundred and eight pairs of strings of 8 bp were obtained, and 26 of them proved to belong to ungapped alignment restricted case (in both string all letters appears at least one time). The number of matches varied from 0 ( $15 \%$ [3.99; 34.47], where the lower and upper bound of $95 \%$ confidence interval calculated using an exact approach (Jäntschi and Bolboacă 2010) are provided in square brackets) to 5 (0.15; 26.78) with the highest frequency at two matches (38\% [19.38; 57.54]). According to full enumeration results, for $n=8$ and $q=4$, the matches are not by chance if more than 3 are observed. More than three matches were observed in four out of twenty-six cases ( $15 \%$ [3.99; 34.47]) on AHSP experiment, leading to the conclusions that these ungapped alignments between HS and MP are not by chance.

Although there is much remains to be done, the work presented in this manuscript generates findings in the field of distribution analysis on equal length sequence alignment. The case with and without restriction were investigated for the number of letters in the alphabet that varied from 2 to 4. Although the present study provides full results for the unrestricted case of ungapped alignment and yielded some finding for the restricted case, its design is not without deficiencies. One of the main limitations of our study is relatively small size of the studies samples ( $n \leq 10$ ) but this is linked with the applied method, full enumeration. Other main limitation is the lack of full characterization of the restricted case. Even if most formulas were identified and can be verified using the results obtained by full enumeration, due to the complexity of the calculations, we did not succeed yet to identify the general formulas for probability mass function and cumulative distribution function for
restricted case, these two statistics being under investigation in our lab. Although the formulas for some alignment statistical parameters were identified (mean and mode) for both investigated cases (unrestricted and restricted case), the variance formula has been identified just for the unrestricted case. Investigation of other statistical parameters such as skewness and kurtosis could also bring valuable information regarding the distribution of alignments. Thus, these statistics could be used as approximation method (e.g. Fisher-Tippet (1928) with the same skewness and kurtosis) of the distribution functions in limit cases. The approximate formulas for the alignment by chance had been obtained for restricted case. Furthermore, it could also be interesting to extend the research regarding the alignment by chance for the particular case of amino-acids sequences. Despite its limitations, this study can be seen as the first step in assessment of the distribution analysis of real ungapped sequences alignments.

## 4 Conclusions

General formula for total number of possibilities of arrangements between two sequences of identical length $\Sigma_{0 \leq i \leq n} S q(i ; n, q)$, total number of matches $S q(i ; n, q)$, and perfect alignment $S q(i=n ; n, q)$ had been identified for both unrestricted case and restricted case, for number of letters in the alphabet from 2 to 4 . Formulas of mean and mode associated to $S q(i ; n, q)$ had been identified for both cases while the formula of variance had been identified just for the unrestricted case. Furthermore, the probability mass function and cumulative distribution function had also been identified for the unrestricted case, while approximate formulas for restricted alignment by chance are presented for $2 \leq q \leq 4$ (where $q=$ number of letters in the alphabet).

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